SQL as Data Manipulation Language (DML)

Insert and update data

Simple SQL queries

Advanced SQL queries

Views

SQL / DML: Overview

- Insert, update, delete data
- Query data
  - Interactively
  - (Embedded in host language)
- Data presentation to users
  - Output improvements
  - Views

SQL / DML: Insert data

- Complete form:
  - Predefined order of values

```
INSERT INTO Customer
VALUES (001, 'Müller', 'Tina', NULL,NULL);
```

- Incomplete form:
  - Free order of values

```
INSERT INTO Customer
(last_name, mem_no) VALUES ('Müller', 001);
```

SQL / DML: Insert data

- Inserting dates

```
INSERT INTO movie
VALUES (95, 'Psycho', 'suspense',
             TO_DATE('1969', 'yyyy'),
            'Hitchcock', 2.00, NULL);
```

- Conversion functions
  - String to date

    `TO_DATE(<string>[,<format>])  ORACLE/Postgres`

  - Date to string

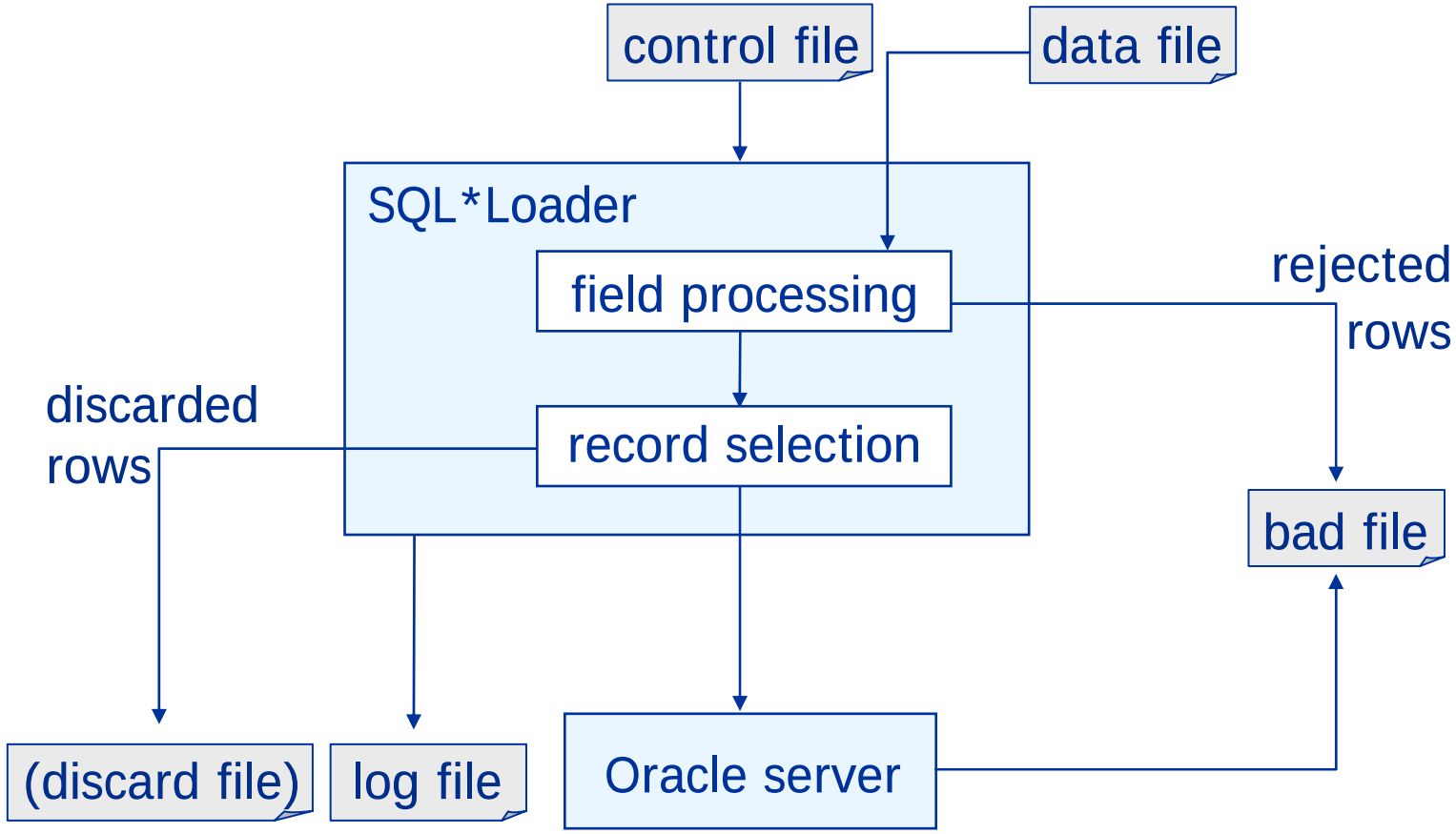
    `TO_CHAR(<date>[,<format>]) ORACLE/Postgres`

SQL / DML: Insert data

- Loading data from files
  - System dependent
- Oracle:
  - **`INSERT INTO …`**
  - Bulk load from file: SQL loader
  - Bulk load from other database: export / import tool
- MySQL:
  - **`INSERT INTO …`**
  - Bulk load: **`LOAD DATA …`**
  - Bulk load from other database:
    `SELECT … INTO OUTFILE, LOAD DATA…`
- Postgres:
  - **`INSERT INTO …`**
  - Bulk load: **`Copy …`**
  - Bulk load from other database: Copy …

SQL / DML: Implementations of bulk load

- Oracle SQL loader



SQL / DML: Implementations of bulk load

- Example:

```
CREATE TABLE loadtest(
    name varchar(20),
    num number(10,2));
```

loadtest.dat

```
'four' , 4
'five' , 5
'six'  , 6
```

- Oracle Syntax:

```
sqlldr <user>/<password> <controlfile>
       <logfile> <badfile> <datafile>
```

loadtest.ctl

```
load data
infile 'loadtest.dat'
badfile 'loadtest.bad'
discardfile 'loadtest.dis'
APPEND INTO table loadtest
 fields terminated by " , "
 optionally enclosed by " ' "
(name char, num integer external)
```

SQL / DML: Implementations of bulk load

- MySQL Example:

```
Mysql> LOAD DATA INFILE 'loadtest.dat'
    ->    IGNORE
    ->    INTO TABLE loadtest
    ->    FIELDS
    ->      TERMINATED BY ","
    ->      OPTIONALLY ENCLOSED BY ''
    ->     (name, num );
Query OK, 3 rows affected (0.01 sec)
Records: 3 Deleted:0 Skipped: 0 Warnings: 0
```

SQL / DML: Insert unique data

- For synthetic keys, e.g. tapeId, …
- Counter-variables in application not sufficient
  - Session dependent
  - Concurrent access problematic
- Counter relation in database
  - Expensive
- Proprietary solutions in existing DBMS
  - MySQL: autoincrement keyword
  - SQL-Server: identity keyword
  - PostgreSQL: serial type and sequence
  - Oracle: sequence-object

SQL / DML: Insert unique data - sequence

- Abstract sequence-object (Oracle)
  - Creates unique integer values

- Syntax:

```
CREATE SEQUENCE <seqName>
[START WITH <integer>]
[INCREMENT BY <integer>]
[MAXVALUE <integer> | NOMINVALUE]
[MINVALUE <integer> | NOMAXVALUE]
[CYCLE | NOCYCLE]
[CACHE <integer> | NOCACHE]
[ORDER | NOORDER];
```

- Example:

```
create sequence tape_sequence;
```

SQL / DML: Insert unique data - sequence

- Value Access
  - **<seqName>.NEXTVAL** = value of last call + increment
  - **<seqName>.CURRVAL** = value of last call

  - **SELECT <seqName>.CURRVAL FROM DUAL;**
  - **DUAL** is Oracle pseudo-table

- Example:

```
INSERT INTO tape
VALUES(tape_sequence.nextval, 'DVD', 95);
```

SQL / DML: Delete data

- Syntax:

```
DELETE from <tableName>
[WHERE <predicate>];
```

  - Delete all rows :

```
DELETE from <tableName>;
```

- Example:

```
DELETE from tape
WHERE format= 'Beta';
```

SQL / DML: Update data

- Syntax:

```
UPDATE <tableName>
SET  <attr> = <value>
   {,<attr> = <value> }
WHERE <predicate>
```

- Examples:

```
UPDATE Customer
SET telephone = 456789
WHERE mem_no = 200;
```

```
UPDATE Rental
SET until_date = SYSDATE
WHERE tape_ID = 3
AND mem_no = 200
AND TO_CHAR(from_date,'yyyy-mm-dd')='2002-05-01';
```

SQL / DML: Example database

```
insert into customer values (001, 'Müller', 'Tina', NULL,NULL);
insert into customer values (007, 'Katz', 'Anna', NULL,NULL);
insert into customer values (002, 'Maus', 'Carla', NULL,NULL);
....

insert into movie values (95, 'Psycho', 'suspense',
            to_date('1969', 'yyyy'), 'Hitchcock', 2.00, NULL);
insert into movie values (112, 'ET', 'comedy',
            to_date('1982', 'yyyy'), 'Spielberg', 1.50, NULL);
....

insert into format values('DVD', '2.00');
insert into format values('Beta', '0.00');
insert into format values('VHS', '0.00');


create sequence tape_sequence;
insert into tape values (tape_sequence.nextval, 'DVD', 95);
insert into tape values (tape_sequence.nextval, 'DVD', 112);
insert into tape values (tape_sequence.nextval, 'VHS', 222);
```

SQL / DML: Example database

```
insert into rental values (3, 1,
                        to_date('2002-05-01','yyyy-mm-dd'), NULL);
insert into rental values (4, 1,
                        to_date('2002-05-01','yyyy-mm-dd'), NULL);
insert into rental values (5, 3,
                        to_date('2002-05-01','yyyy-mm-dd'),
                        to_date('2002-05-02','yyyy-mm-dd'));


insert into actor values ('Hitchcock',
                    'Hitchcock', to_date(1899-08-13','yyyy-mm-dd'));
insert into actor values ('Harrison Ford',
                'Harrison Ford', to_date('1942-07-13','yyyy-mm-dd'));


insert into play values(290,'Harrison Ford');
insert into play values(98,'Hitchcock');
```

SQL / DML: Querying Language

- SQL is relational complete
- Additional query concepts
  - Advanced search expressions on strings
    e.g., find all movies starting with "star wars"
  - Arithmetic in expressions,
    e.g., number of tapes for each movie
  - Grouping and predicates over sets
    e.g., total receipts of each movie within the last year

SQL / DML: Basics

- Basic query pattern:

```
SELECT  [DISTINCT] A1, A2,...,An
FROM    R1, R2,...Rm
WHERE   predicate P;
```

  - A_1 , A_2 , ..., A_n attribute names,
  - R_1 , R_2 , ..., R_m relation names,
  - P Boolean predicate on attributes and constants

- Equivalent to relational algebra expression:

$$\Pi_{A1, A2, ..., An} (\sigma_P (R_1 \times R_2 \times ... \times R_m))$$

  - Projection (RA) → **SELECT** (SQL)
  - Cartesian Product (RA) → **FROM** (SQL)
  - Selection (RA) → **WHERE** (SQL)

SQL / DML: Basics

- Query result is relation
- Query evaluation order:
  1. **FROM**-clause
  2. **WHERE**-clause
  3. **SELECT**-clause
- No duplicate removal (performance!)

```
SQL> SELECT last_name
  2   FROM Customer;

LAST_NAME
------------------------------
Müller
Katz
Maus
Hinz
Kunz
Müller
```

SQL / DML: Basics

- Eliminating duplicates:
  - Targetlist contains KEY attribute
  - Targetlist constrains UNIQUE attribute
  - Targetliste defined with **DISTINCT**

```
SELECT mem_no, last_name
FROM Customer
```

```
 MEM_NO     LAST_NAME
---------- ----------------
         1 Müller
         7 Katz
         2 Maus
        11 Hinz
        23 Kunz
       111 Müller
```

```
SELECT DISTINCT last_name
FROM Customer
```

```
LAST_NAME
----------------
Hinz
Katz
Kunz
Maus
Müller
```

SQL / DML: Basics

- **WHERE**-clause structure:
  - Simple Boolean predicates similar to RA and Calculus
  - Additional simple predicates:

    <attribute> BETWEEN <value1> AND <value2>
    <attribute> IS [NOT] NULL
    <attribute> LIKE <string>
    <attribute> SIMILAR TO <string>

  - Advanced predicated with sub-queries

    Set-operators (**IN, NOT IN, SOME, ALL, EXISTS**)

SQL / DML: Simple queries

Example: All customers named Anna

```
SQL>  select mem_no, last_name, first_name
  2   from customer
  3   where first_name='Anna';

    MEM_NO LAST_NAME                      FIRST_NAME
---------- ------------------------------ --------------
         7 Katz                           Anna
        23 Kunz                           Anna
```

Example: All movies by Lucas from 1999 or later

```
SQL> select id, title
  2  from movie
  3  where director='Lucas'
  4  and to_char(year,'yyyy')>='1999';

        ID TITLE
---------- ----------------------------------------
       345 Star Wars I
```

SQL / DML: Simple queries

- More examples:

All formats with extra charge between 1 and 2 Euro

```
SELECT *
FROM Format
WHERE charge BETWEEN 1.00 and 2.00;
```

All tapes currently on loan

```
SELECT tape_id
FROM Rental
WHERE until_date IS NULL;
```

SQL / DML: Simple queries - expressions

Core SQL:1999

- **LIKE** - expression
  - Simple form of regular expression
  - % : any sequence of characters
  - _ : exactly one character

- Example:

All 'star wars' movies

```
SQL> select id, title, director
  2  from movie
  3  where title like 'Star Wars %';

        ID TITLE        DIRECTOR
---------- ------------ ---------------
       345 Star Wars I  Lucas
       290 Star Wars IV Lucas
```

SQL / DML: Simple queries - expressions

enhanced SQL:1999

- **SIMILAR** - expression
  - Advanced form of regular expression

- Example:

  All 'star wars' movies

```
SELECT id, title, director
FROM movie
WHERE title SIMILAR TO
      'Star Wars (I | IV | V | VI | 1 | [4-6])';
```

SQL / DML: Simple queries

- Member in set: **IN**

All movies from Spielberg or Lukas

Core SQL:1999

```
SELECT title, director
FROM Movie
WHERE director IN ('Spielberg','Lucas');
```

All movies from Lucas in 1999

enhanced SQL:1999

```
SELECT title, director
FROM Movie
WHERE (director, year)
  IN (('Lucas',to_date(1999,'yyyy')));
```

SQL / DML: Simple queries - expressions

- Functions
  - Expressions may contain functions
  - Arithmetical and string built-in functions
  - User defined functions on user defined types

- String function examples:
  - **SOUNDEX (<string>), UPPER (<string>)**
  - **SUBSTRING(<string> FROM <integer> FOR <integer>)**

```
SQL> SELECT title, director
  2  FROM Movie
  3  WHERE SOUNDEX(director)  =  SOUNDEX('Spilbak');

TITLE              DIRECTOR
------------------ ----------------
ET                 Spielberg
Psycho             Spielberg
Jaws               Spielberg
```

SQL / DML: Simple queries - expressions

- Arithmetic function examples
  - **SQRT(<number>)**
  - Basic arithmetic expressions

All tapes, their price and tax

```
SQL> SELECT id, pricepday,
  2         0.16*pricepday as tax
  3  FROM Movie;

        ID  PRICEPDAY        TAX
---------- ---------- ----------
        95          2        .32
       112        1.5        .24
       345          2        .32
       222        2.2       .352
       290          2        .32
       100        1.5        .24
```

SQL / DML: Simple queries - expressions

- Date function examples
  - differ heavily between systems
  - Oracle: **SYSDATE, MONTHS_BETWEEN, ADD_MONTHS**

```
SQL> SELECT title, to_char(year,'yyyy') as year
  2  FROM movie
  3  WHERE months_BETWEEN(SYSDATE,year)> 120 ;

TITLE                     YEAR
------------------------- ----
Psycho                    1969
ET                        1982
Jaws                      1975
```

SQL / DML: Simple queries

Core SQL:1999

- Combination of relations
  - Schema compatible relations
  - **UNION, INTERSECT, EXCEPT**

- Syntax:

```
UNION | INTERSECT | EXCEPT
          [DISTINCT | ALL]
          [CORRESPONDING [BY <attributes>]]
```

  - Default **DISTINCT**

  - Default: all attributes used
  - **CORRESPONDING BY**: defines used common attributes
  - **CORRESPONDING**: uses all common attributes

SQL / DML: Simple queries

- Example:

All movies from Spielberg or Lukas

```
(SELECT title, director
 FROM Movie
 WHERE director like 'Spielberg')
UNION
(SELECT title, director
 FROM Movie
 WHERE director like 'Lucas');
```

SQL / DML: Simple queries

- More examples:

  All movies not by Lucas

```
(SELECT *
 FROM Movie)
EXCEPT
(SELECT * from Movie
 WHERE director='Lucas');
```

  All directors and actors in our database

```
(SELECT director as celebrity
 FROM Movie)
UNION DISTINCT
(SELECT stage_name as celebrity
 FROM Actor);
```

SQL / DML: Implementations combinations

- Oracle:
  - **UNION**
  - **MINUS** implements **EXCEPT**
  - **INTERSECT**
  - **CORRESPONDING [BY]** not implemented
- PostgreSQL:
  - **UNION**
  - **EXCEPT**
  - **INTERSECT**
  - **CORRESPONDING [BY]** not implemented
- MySQL:
  - **UNION, EXCEPT, INTERSECT** not implemented

SQL / DML: Simple queries with joins

Core SQL:1999

- Simple joins
  - Search predicates and join conditions mixed

Example: All Tapes and their corresponding movie

```
SQL> SELECT t.id, t.format, m.id, m.title
  2  FROM Tape t, Movie m
  3  WHERE m.id = t.movie_id;

        ID FORMAT         ID TITLE
---------- ----- ---------- ----------------------------
         1 DVD            95 Psycho
         2 DVD           112 ET
         3 VHS           222 Psycho
         4 DVD           345 Star Wars I
         5 VHS           345 Star Wars I
         9 VHS           345 Star Wars I
```

SQL / DML: Simple queries with joins

enhanced SQL:1999

- Cross join (cross product)

```
<tableName> CROSS JOIN <tableName>
```

- Natural inner join

```
<tableName> NATURAL [INNER] JOIN <tableName>
```

- Example:

```
SELECT *
FROM Rental NATURAL INNER JOIN Customer;
```

SQL / DML: Simple queries with joins

enhanced SQL:1999

- Inner join with attribute list

```
<tableName> [INNER] JOIN <tableName>
USING <attributList>
```

Subset of attributes in common

- Example:

```
SELECT *
FROM Rental r JOIN Customer c
USING (mem_no);
```

SQL / DML: Simple queries with joins

enhanced SQL:1999

- Inner join with condition

```
<tableName> [INNER] JOIN <tableName>
ON <condition>
```

- Examples:

```
SELECT *
FROM Tape t JOIN Movie m
ON t.movie_id = m.id;
```

```
SELECT *
FROM Rental r JOIN Customer c
ON r.mem_no = c.mem_no;
```

SQL / DML: Simple queries with joins

All Customers who have rented at least one science fiction film

```
SELECT c.mem_no, c.last_name, c.first_name
FROM ((Customer c
          JOIN Rental r ON c.mem_no = r.mem_no)
          JOIN Tape t ON t.id = r.tape_id )
          JOIN Movie m ON t.movie_id = m.id
WHERE m.category='Scifi';
```

```
SELECT c.mem_no, c.last_name, c.first_name
FROM Customer c, Rental r, Tape t, Movie m
WHERE c.mem_no=r.mem_no
AND t.id = r.tape_id
AND t.movie_id = m.id
AND m.category='Scifi';
```

SQL / DML: Simple queries with joins

enhanced SQL:1999

- Natural outer join

```
<tableName> LEFT|RIGHT|FULL
NATURAL [OUTER] JOIN <tableName>
```

- Outer join with condition

```
<tableName> LEFT|RIGHT|FULL [OUTER] JOIN <tableName>
ON <condition>
```

- Example:

```
SELECT *
FROM Rental r RIGHT OUTER JOIN Customer c
ON r.mem_no = c.mem_no;
```

SQL / DML: Simple queries with joins

- Example (extended)

```
SQL> SELECT r.tape_id, r.from_date, c.mem_no,c.first_name
  2  FROM Rental r RIGHT OUTER JOIN Customer c
  3  ON r.mem_no = c.mem_no;

   TAPE_ID FROM_DATE     MEM_NO FIRST_NAME
---------- --------- ---------- --------------------
         3 01-MAY-02          1 Tina
         4 01-MAY-02          1 Tina
         5 01-MAY-02          2 Carla
                             23 Anna
                            111 Bert
                             11 Fritz
                              7 Anna
```

SQL / DML: Implementations of joins

- Oracle/Postgres:
  - Simple join
  - Cross join
  - (natural) inner join with attribute list, with condition
  - (natural) Right, left, full outer join with condition

  - recommends ANSI-syntax for compatibility

- MySQL:
  - Simple join
  - Cross join
  - Straight join (left table always read before right one)
  - Inner join with condition
  - (natural) left, right outer join with condition

SQL / DML: Improving the output

- Not feature of relational algebra
- Example:
  - Rename column title for *this* query
  - Order tuples

```
SQL> SELECT m.title as Movies, t.id, t.format
  2  FROM Movie m, Tape t
  3  WHERE m.id = t.movie_id
  4  ORDER BY title;

MOVIES                                        ID FORMAT
---------------------------------------- -------- ------
ET                                             2 DVD
Psycho                                         1 DVD
Psycho                                         3 VHS
Star Wars I                                    4 DVD
Star Wars I                                    5 VHS
Star Wars I                                    9 VHS
```

SQL / DML: Improving the output

- Syntax:

```
ORDER BY <orderexpression> ASC|DESC
```

- Ordering expression
  - No advanced expressions (no sub-query, no grouping)
  - At least one attribute reference
  - References in order expression $\subseteq$ result attributes

- Multiple sort attributes:
  - Primary ordering by first attribute
  - Secondary ordering by second attribute, …

```
SELECT m.title as Movies, t.id, t.format
FROM Movie m, Tape t
WHERE m.id = t.movie_id
ORDER BY title, format;
```

SQL / DML: Improving the output

- Advanced features system dependent
- Oracle SQL+ Example:
  - Format column title for *all* queries
  - Don't repeat identical titles

```
SQL> BREAK ON title
SQL> COLUMN title HEADING "Movies" FORMAT A15

SQL> SELECT m.title, t.id, t.format
  2  FROM Movie m, Tape t
  3  WHERE m.id = t.movie_id;

Movies                      ID FORMAT
--------------- ---------- ------
Psycho                       1 DVD
ET                           2 DVD
Psycho                       3 VHS
Star Wars I                  4 DVD
                             5 VHS
                             9 VHS
```

SQL / DML: Sub-queries

Core SQL:1999

- Sub-queries with single results
  - Operators {=, ≤, ≥, ≠, <, >}
  - Expressible without sub-query

- Example:

  Movies shorter than 'Star Wars I' (id 345)

```
SELECT m.id
FROM Movie m
WHERE m.length <
                (SELECT m1.length
                 FROM Movie m1
                 WHERE m1.id = 345);
```

SQL / DML: Sub-queries

Core SQL:1999

- Set Operator **IN**

- Independent sub-query example:

All Tapes for movie 'Star Wars'

```
SELECT t.id, t.format
FROM Tape t
WHERE t.movie_id
      IN (SELECT m.id
          FROM Movie m
          WHERE m.title like 'Star Wars %');
```

SQL / DML: Sub-queries

Core SQL:1999

- Set Operator **IN**

- Correlated sub-query example:

  Directors playing in their own movies

```
SELECT m.director
FROM Movie m
WHERE m.director IN
      (SELECT p.actor_name
       FROM Play p
       WHERE p.movie_id = m.id);
```

SQL / DML: Sub-queries

Core SQL:1999

- Alternative syntax: **EXISTS**
- Example:

```
SELECT t.id, t.format
FROM Tape t
WHERE EXISTS (SELECT *
              FROM Movie m
              WHERE t.movie_id = m.id
              AND m.title like 'Star Wars %');


SELECT m.director
FROM Movie m
WHERE EXISTS (SELECT *
              FROM Play p
              WHERE p.movie_id = m.id
              AND m.director like p.actor_name);
```

SQL / DML: Query rewriting

- Rewriting possible for **IN, EXISTS**

- Examples:

```
SELECT t.id, t.format
FROM Tape t, Movie m
WHERE m.id = t.movie_id
AND m.title like 'Star Wars %';
```

```
SELECT m.director
FROM Movie m, Play p
WHERE p.movie_id = m.id
AND m.director like p.actor_name;
```

SQL / DML: Sub-queries

- Negation **NOT EXISTS, NOT IN**

Core SQL:1999

All tapes that never have been on loan

```
SELECT t.id
FROM Tape t
WHERE NOT EXISTS (SELECT *
                  FROM Rental r
                  WHERE r.tape_id = t.id);
```

All movies no copy of which are currently on loan

```
SELECT distinct m.id
FROM Movie m
Where NOT EXISTS (SELECT *
                  FROM Rental r, Tape t
                  WHERE r.tape_id = t.id
                  AND m.id = t.movie_id
                  AND r.until_date IS NULL);
```

SQL / DML: Quantified sub-queries

enhanced SQL:1999

- Quantified comparison operators
  - No re-writing without sub-query
  - Quantification: **ALL**, **SOME** (synonym **ANY**)
  - Operators $\in \{=, \leq, \geq, \neq, <, >\}$

- Quantification **All**
  - Sub-query true if true for all tuples

Most expensive movies

```
SELECT m.id, m.pricepday
FROM Movie m
WHERE m.pricepday >= ALL
                     (SELECT m1.pricepday
                      FROM MOVIE m1);
```

SQL / DML: Quantified sub-queries

enhanced SQL:1999

- Quantification: **SOME**
  - Sub-query true if true for at least one tuple

```
SELECT title, pricePDay
FROM Movie m
WHERE pricePday  < SOME
                   (SELECT m1.pricepday
                    FROM Movie m1);
```

  - **= SOME** *equivalent* **IN**

```
SELECT m.director
FROM Movie m
WHERE m.director = SOME
      (SELECT p.actor_name
       FROM Play p
       WHERE p.movie_id = m.id);
```

SQL / DML: Implementations of sub-queries

- Oracle:
  - Sub-queries with single results
  - Sub-queries with **[NOT] IN, [NOT] EXISTS**
  - Quantified comparison **ALL, SOME, ANY**
- PostgreSQL
  - Similar to Oracle
- MySQL:
  - No sub-queries supported

SQL / DML: Universal quantifiers

- ∀ - Quantification
- ∃ - Quantification (exactly one)
  - Describe counterexample
  - Combine with **NOT EXISTS**

Movies with only one tape

```
SELECT m.id
FROM Movie m
WHERE NOT EXISTS
   (SELECT *
    FROM Tape t1, Tape t2
    WHERE t1.movie_id = t2.movie_id
    AND t1.id <> t2.id
     AND t2.movie_id = m.id );
```

SQL / DML: Universal quantifiers

All Customers whose rented movies all have category "suspense"

```
SELECT c.mem_no
FROM Customer c
WHERE NOT EXISTS
          (SELECT m.id
           FROM Movie m, Rental r, Tape t
           WHERE m.id = t.movie_id
           AND r.tape_id = t.id
           AND c.mem_no = r.mem_no
           AND m.category <> 'suspense');
```

SQL / DML: Universal quantifiers

Customers that had rented all movies

```
SELECT c.mem_no
FROM Customer c
WHERE NOT EXISTS
          (SELECT m.id
           FROM Movie m
           WHERE NOT EXISTS
                    (SELECT *
                     FROM Rental r, Tape t
                     WHERE m.id = t.movie_id
                     AND r.tape_id = t.id
                     AND c.mem_no = r.mem_no));
```

SQL / DML: Universal quantifiers

Customers that rented only one movie

```
SELECT c.mem_no
FROM Customer c, Rental r, Tape t, Movie m
WHERE c.mem_no = r.mem_no
AND r.tape_id = t.id
AND t.movie_id = m.id
AND NOT EXISTS
          (SELECT m1.id
           FROM Rental r1, Tape t1, Movie m1
           WHERE r1.tape_id = t1.id
           AND t1.movie_id = m1.id
           AND c.mem_no = r1.mem_no
           AND m1.id <> m.id);
```

SQL / DML: Aggregate functions

- Mathematical aggregate functions on data sets
- Example: SUM, AVG, MIN, MAX, COUNT
- Not in relational algebra

```
SQL> SELECT MIN(pricePDay) as MIN,
  2  MAX(pricePDay) as MAX, AVG(pricePDay)
  3  FROM Movie;

       MIN        MAX AVG(PRICEPDAY)
---------- ---------- --------------
       1.5        2.2     1.86666667
```

- Target list: only aggregate functions or none
  - Exception: **GROUP BY**

SQL / DML: Aggregate functions

- Comparison using aggregates: sub-queries

Movies with price above average

```
SQL> SELECT m.id, m.Title, m.pricepday
  2  FROM Movie m
  3  WHERE pricePDay >
  4                   (SELECT AVG(pricePDay)
  5                    FROM Movie);

     ID TITLE                                PRICEPDAY
------- ------------------------------------ ---------
     95 Psycho                                       2
    345 Star Wars I                                  2
    222 Psycho                                     2.2
    290 Star Wars IV                                 2
```

SQL / DML: Aggregate functions

- Examples:

Movies with minimal price

```
SELECT m.Title, m.pricepday
FROM Movie m
WHERE pricePDay =
                    (SELECT MIN(pricePDay)
                     FROM Movie);
```

Movie with more than 2 tapes

```
SELECT m.id, m.title
FROM Movie m
WHERE 2 <  (SELECT count(t.id)
            FROM tape t
            WHERE t.movie_id = m.id)
```

SQL / DML: Aggregate functions

- More examples:

Movies having tapes in one format only

```
SELECT m.id, m.title
FROM Movie m, Tape t1
WHERE m.id = t1.movie_id
AND 0 =
        (SELECT COUNT(*)
         FROM Tape t2
         WHERE t1.id <> t2.id
         AND t1.format <> t2.format
         AND t2.movie_id = m.id);
```

SQL / DML: Aggregate functions

- Additional qualification with **DISTINCT | ALL**

- Example:

  Movies that are available in all formats

```
SELECT DISTINCT t1.movie_id
FROM Tape t1
WHERE
     (SELECT COUNT(DISTINCT format)
      FROM   Tape t2
      WHERE  t2.movie_id = t1.movie_id)
      =
      (SELECT COUNT(*)
       FROM Format);
```

SQL / DML: Grouping

- Syntax:

```
SELECT <targetlist>
FROM <tablelist>
[WHERE <predicate>]
GROUP BY <attributelist>
```

- Groups all rows with same values in <attributelist>
- Target list: grouping attributes and aggregates

- Example:

  Number of tapes in each format

```
SELECT t.format, count(t.id)
FROM Tape t
GROUP BY t.format;
```

SQL / DML: Grouping

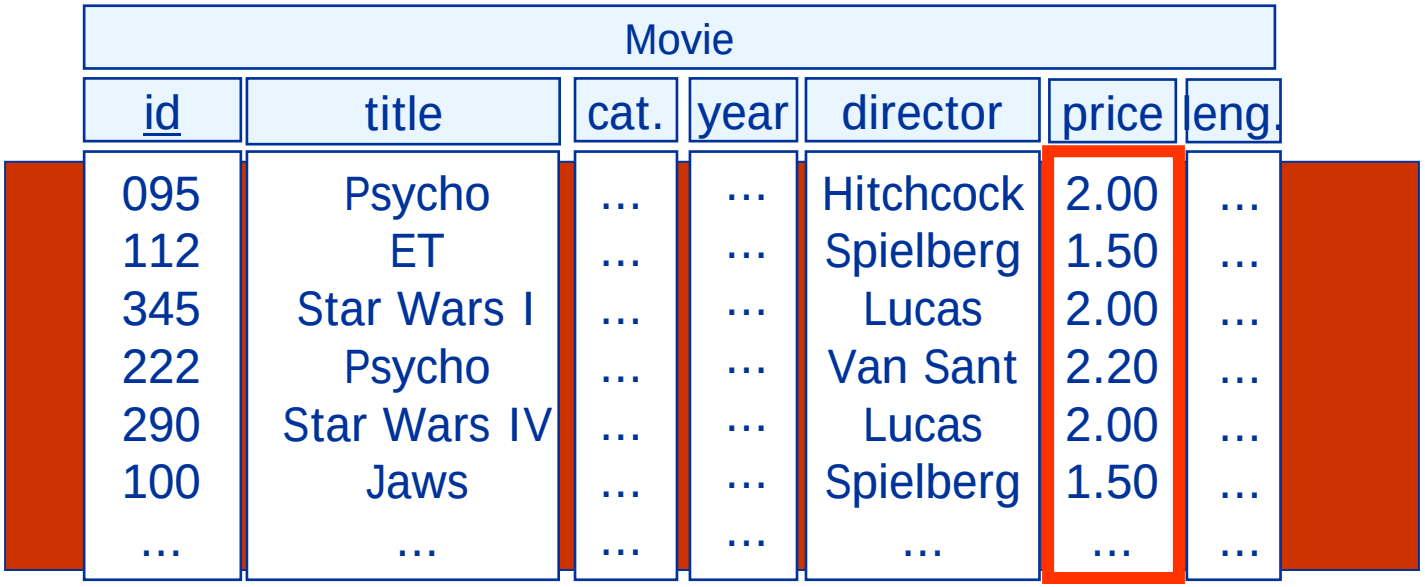
- Aggregates evaluated over groups

Number of tapes for each movie

```
SQL> SELECT t.movie_id, count(*)
  2  FROM Tape t
  3  GROUP BY t.movie_id;

  MOVIE_ID   COUNT(*)
---------- ----------
        95          1
       100          1
       112          1
       222          1
       290          1
       345          4
```

SQL / DML: Grouping



Movie

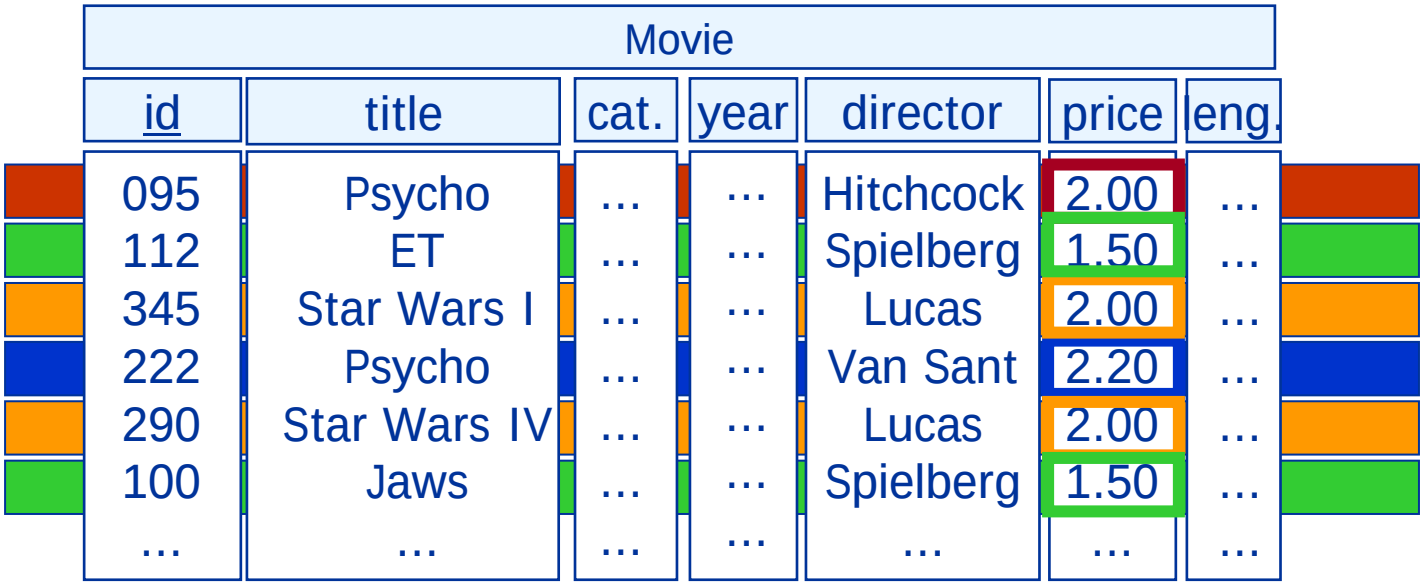
id	title	cat.	year	director	price	leng.
095	Psycho	...	...	Hitchcock	2.00	...
112	ET	...	...	Spielberg	1.50	...
345	Star Wars I	...	...	Lucas	2.00	...
222	Psycho	...	...	Van Sant	2.20	...
290	Star Wars IV	...	...	Lucas	2.00	...
100	Jaws	...	...	Spielberg	1.50	...
...	...	...	...	...	...	...

```
SELECT sum(price)
FROM movie;
```

11.2

Implicit group: all tuples in table

SQL / DML: Grouping



Movie

id	title	cat.	year	director	price	leng.
095	Psycho	...	...	Hitchcock	2.00	...
112	ET	...	...	Spielberg	1.50	...
345	Star Wars I	...	...	Lucas	2.00	...
222	Psycho	...	...	Van Sant	2.20	...
290	Star Wars IV	...	...	Lucas	2.00	...
100	Jaws	...	...	Spielberg	1.50	...
...	...	...	...	...	...	...

```
SELECT director, sum(price)
FROM movie
Group by director;
```

Hitchcock	2.0
Spielberg	3.0
Lucas	4.0
Van Sant	2.2

SQL / DML: Grouping

Total receipts of each tape within the last year

```
SQL> SELECT t.id, count(*)
  2  FROM Tape t, Rental r
  3  WHERE t.id = r.tape_id
  4  AND to_char(r.from_date,'yyyy') >= 2005
  5  GROUP BY t.id;

        ID   COUNT(*)
---------- ----------
         1          1
         2          1
         3          2
         4          2
         5          1
        11          1
        12          1
```

SQL / DML: Grouping

Total receipts of each movie within the last year

```
SQL> SELECT t.movie_id, count(*)
  2  FROM Tape t, Rental r
  3  WHERE t.id = r.tape_id
  4  AND to_char(r.from_date,'yyyy') >= 2005
  5  GROUP BY t.movie_id;

  MOVIE_ID   COUNT(*)
---------- ----------
        95          1
       100          1
       112          1
       222          2
       290          1
       345          3
```

SQL / DML: Grouping + Having

- Qualifying predicate for groups

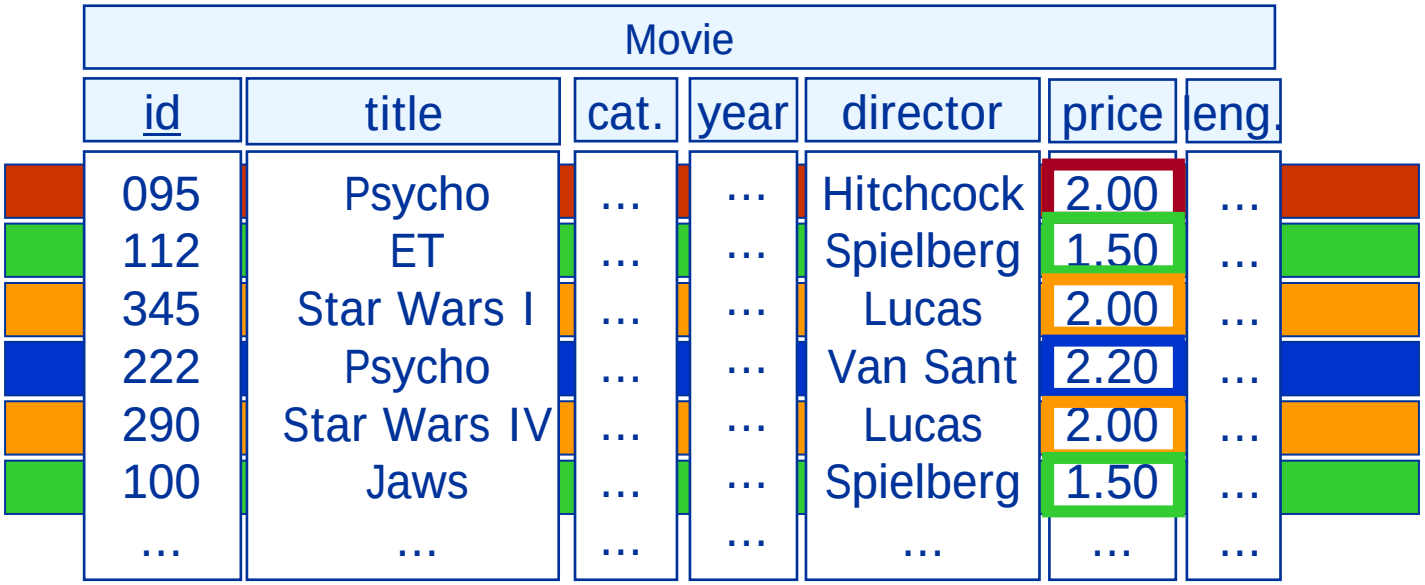
```
SQL> SELECT f.name, sum(charge)
  2    FROM Rental r, Tape t,  format f
  3    WHERE t.id = r.tape_id
  4    AND t.format=f.name
  5    GROUP BY f.name;

NAME  SUM(CHA
----- ------
Beta
DVD
VHS
```

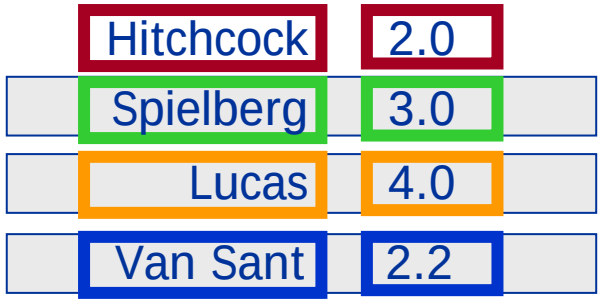
```
SQL> SELECT f.name, sum(charge)
  2    FROM Rental r, Tape t,  format f
  3    WHERE t.id = r.tape_id
  4    AND t.format=f.name
  5    GROUP BY f.name
  6    having count(f.name)>2;

NAME  SUM(CHARGE)
----- -----------
DVD             8
VHS             0
```

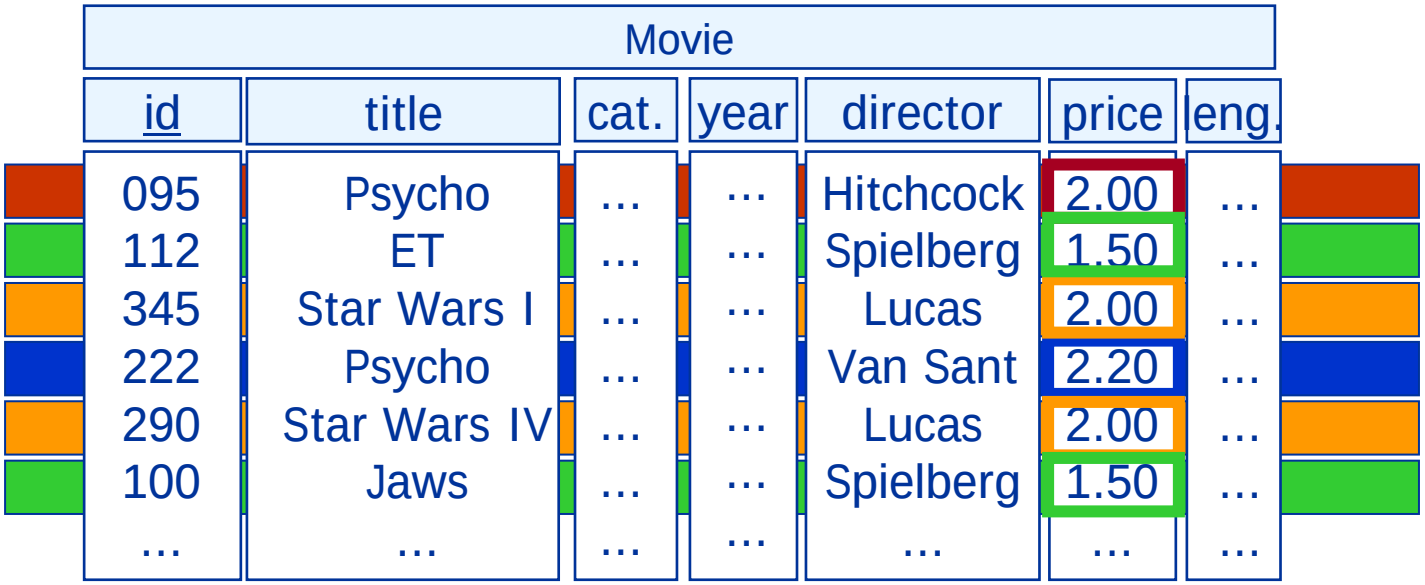
SQL / DML: Grouping + Having



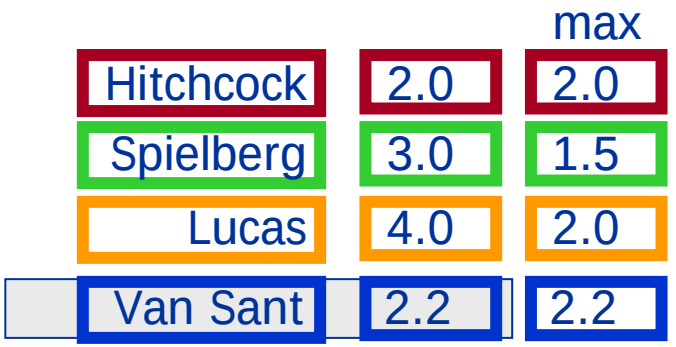

```
SELECT director, sum(price)
FROM movie
Group by director
HAVING sum(price)>2.00;
```



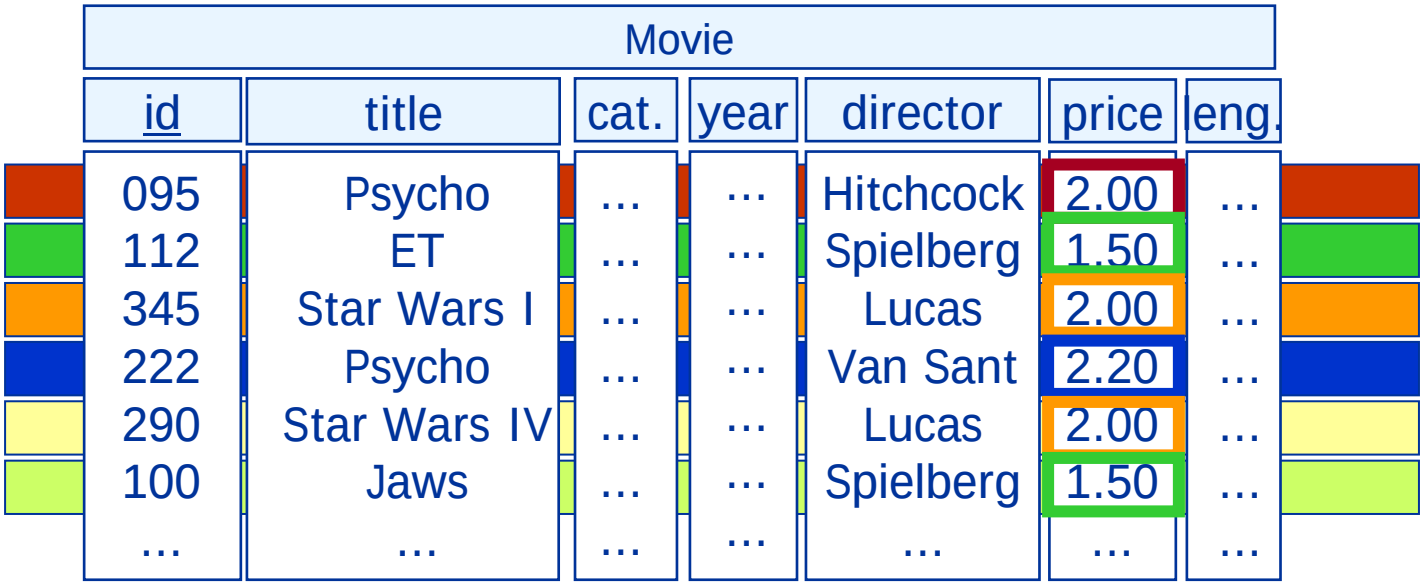
SQL / DML: Grouping + Having




```
SELECT director, sum(price)
FROM movie
Group by director
HAVING max(price)>2.00;
```



SQL / DML: Grouping + Having



Movie

id	title	cat.	year	director	price	leng.
095	Psycho	…	…	Hitchcock	2.00	…
112	ET	…	…	Spielberg	1.50	…
345	Star Wars I	…	…	Lucas	2.00	…
222	Psycho	…	…	Van Sant	2.20	…
290	Star Wars IV	…	…	Lucas	2.00	…
100	Jaws	…	…	Spielberg	1.50	…
…	…	…	…	…	…	…

```
SELECT director,title, sum(price)
FROM movie
Group by director, title
HAVING max(price)>2.00;
```

		max
Hitchcock, Psycho	2.0	2.0
Spielberg, ET	1.5	1.5
Lucas, SW I	2.0	2.0
Van Sant, Psycho	2.2	2.2
Lucas, SW IV	2.0	2.0
Spielberg, Jaws	1.5	1.5

SQL / DML: Grouping + Having

- **HAVING** without **GROUP BY**
  - Implicit single group contains all tuples

```
SQL> SELECT sum(charge)
  2    FROM Rental r, Tape t,  format f
  3    WHERE t.id = r.tape_id
  4    AND t.format=f.name
  5    having count(f.name)>2;

SUM(CHARGE)
-----------
          8
```

SQL / DML: Grouping + Having

- Query evaluation order:
  1. **FROM**-clause
  2. **WHERE**-clause
  3. **GROUP BY**-clause
  4. **HAVING**-clause
  5. **SELECT**-clause

Number of rentals for all customers named Anna or Tina, which rented some tapes more than once

```
SELECT c.mem_no, count(*)
FROM Rental r, Customer c
WHERE r.mem_no= c.mem_no
AND c.first_name ='Anna'
OR c.first_name='Tina'
GROUP BY c.mem_no
HAVING count(DISTINCT r.tape_id)<count(*);
```

SQL / DML: Nested aggregation with groups

- Nested aggregation using groups

Most loaned movie

```
SELECT t.movie_id, count(t.movie_id)
FROM Rental r, Tape t
WHERE r.tape_id = t.id
GROUP BY t.movie_id
HAVING COUNT(t.movie_id) > = ALL
                    (SELECT count(t1.movie_id)
                     FROM Rental r1, Tape t1
                     WHERE r1.tape_id = t1.id
                     Group BY t1.movie_id);
```

SQL / DML: Nested aggregation with groups

Movie with maximal number of tapes, show number of tapes

```
SELECT m.id, m.title, t1.t_no
FROM (SELECT t.movie_id, count(*) as t_no
      FROM tape t
      GROUP BY t.movie_id) t1, movie m
WHERE m.id=t1.movie_id
AND t1.t_no = (SELECT max(count(*))
               FROM tape t
               GROUP by t.movie_id);
```

SQL / DML: Output improvement

Core SQL:1999

- Select values depending an condition
- Complete **CASE** form:

```
CASE
  WHEN <condition1> THEN <result1>
  [ WHEN <condition2> THEN <result2>
  [ WHEN <condition3> THEN <result3> ]]
  [ ELSE <elseresult>]
END
```

- Example:

```
SELECT length,
CASE WHEN length is NULL then 'not defined'
     WHEN length < 90 THEN 'short'
     ELSE 'long'
END
FROM Movie;
```

SQL / DML: Output improvement

Core SQL:1999

- Simple **CASE** form

```
CASE <operand>
  WHEN <value1> THEN <result1>
  [ WHEN <value2>THEN <result2>
  [ WHEN <value3> THEN <result3> ]]
  [ ELSE <elseresult>]
END
```

- Example:

```
select f.name,
case f.name
when 'DVD' then 'DISC'
when 'Beta' then 'TAPE'
when 'VHS' then 'TAPE'
else NULL
end
from Format f;
```

SQL / DML: Transitive closure

enhanced SQL:1999

- Recursive queries
  - Name recursion expression
  - Use name in associated query expression

- SYNTAX:

```
WITH RECURSIVE
  <queryname1> AS <query1>[,
  <queryname2> AS <query2>,…]
SELECT ...
FROM <queryname1>[,<queryname2>...]
WHERE...
```

SQL / DML: Transitive closure

- Example: All lectures required for lecture XYZ

enhanced SQL:1999

```
create table lecture(
lnr integer primary key,
name varchar(20));

        create table requires(
        pre integer references lecture(lnr),
        suc integer references lecture(lnr),
        constraint req_pk primary key(pre, suc));


WITH RECURSIVE preLecture(pre, suc)
   AS (SELECT pre,suc  FROM requires)
SELECT l.lnr as prerequisite
FROM preLecture p1, preLecture p2, lecture l
WHERE p2.suc = l.lnr
AND   l.name = 'databases'
AND   p1.suc = p2.pre;
```

SQL / DML: Transitive closure

- Different implementation in oracle:

Lecture:

LNR	NAME
1	databases
2	algorithms I
3	algorithms II

Requires:

PRE	SUC
2	3
3	1

```
SQL> SELECT  r.pre
  2  FROM    requires r, lecture l
  3  WHERE   l.name= 'databases'
  4  START WITH r.suc = l.lnr
  5  CONNECT BY PRIOR pre = suc;

       PRE
----------
         3
         2
```

SQL / DML: Structuring

enhanced SQL:1999

- Difficult structuring of complex queries
- No naming of commands / relations for re-use

- Temporary table

```
CREATE TABLE <tablename>
    {global temporary | local temporary }
    <table structure>
[ON COMMIT {PRESERVE ROWS | DELETE ROWS}]
```

  - Stores temporal query result
  - **LOCAL**: Only visible to owner
  - Dropped at end of session

SQL / DML: Temporary table

- Example:

	Test1:
`CREATE TABLE testsource(id integer);` `CREATE GLOBAL TEMPORARY TABLE test1(id integer) ON COMMIT PRESERVE ROWS;`	No rows
`INSERT INTO testsource values(1);`	No rows
`INSERT INTO test1 SELECT * FROM testsource;`	1
`INSERT INTO testsource values(2);`	1
`COMMIT;`	1

SQL / DML: Temporary table

- Example:

	Test2:

`CREATE TABLE testsource(id integer);`
`CREATE GLOBAL TEMPORARY TABLE test2(id integer)`
`ON COMMIT DELETE ROWS;`
	No rows

`INSERT INTO testsource values(1);`	No rows

`INSERT INTO test2`
`SELECT *`
`FROM testsource;`
	1

`INSERT INTO testsource values(2);`	1
`COMMIT;`	No rows

SQL / DML: View

Important concept

- SQL-object (virtual relation)
- Important for
  - Tailoring database schema for different applications
  - Access protection, Privacy
  - Structuring complex queries
- Relational concept for external schema

- Syntax:

```
CREATE VIEW <viewname> AS <query>;
```

- Example:

```
CREATE VIEW rental_overview AS
  SELECT r.mem_no, c.last_name, r.tape_id
  FROM rental r, customer c
  WHERE r.mem_no = c.mem_no;
```

SQL / DML: View updates

Core SQL:1999

- Updateable views:
  1. No **SELECT DISTINCT**
  2. No duplicate attribute in target list
  3. Only one table reference
  4. No **GROUP BY**
  5. No aggregates
  6. No set operators (**INTERSECT, EXCEPT, UNION**)

- Formal:

$$u\,(V(B)) \;=\; V\,(c_u\,(B)\,)$$

  - V(B) view on base tables B
  - u(V(B)) update on view
  - c_u : equivalent update(s) on base relations must exist

SQL / DML: View updates

Core SQL:1999

- Examples:
  - Not updateable (distinct, group by):

```
CREATE VIEW movieformats (movie, numFormats)
  AS SELECT  movie_id, count(distinct format)
  FROM  Tape t
  GROUP BY movie_id;
```

  - Updateable:

```
CREATE VIEW movies (movie, name)
  AS SELECT  id, title
  FROM Movie
  WHERE id > 100;
```

u: `INSERT INTO movies VALUES (47,'The Fugitive');`

c_u:
```
INSERT INTO movie (id, title)
                  VALUES (47,'The Fugitive');
```

SQL / DML: View updates

- Additional conditions:
  - If u does not have an effect, then c_u should not
  - No side effects: c should only effect tuples in B which are represented in V
  - Inverse update: For a view update u there should be an inverse update w such that w (u (V(B))) = V (B)
  - No constraint on base tables must be violated by u

SQL / DML: View updates

- Views with check option
  - CHECK OPTION prevent side effects on base tables
- Syntax:

```
CREATE VIEW <viewname>
AS <query>
WITH CHECK OPTION;
```

- Theoretically more views updateable
  - e.g. UNIQUE columns in joins
- Enhanced SQL:
  - 1999 additional complex conditions for view update

SQL / DML: Remarks about NULL

- NULL treated as "unknown"
- Predicates:
  - NULL AND TRUE = NULL
  - NULL OR TRUE = TRUE
  - predicate evaluates to NULL for row r → r not returned
- Arithmetical expression:
  - If NULL involved → expression evaluates to NULL
- Aggregates:
  - Count(*) counts also NULL
  - avg(*) ≠ sum(*)/count(*)

SQL / DML: Summary

- SQL as Data manipulation language
  - Declarative language
  - Relational complete query language

- Important terms and concepts:
  - Insert, update, delete data
  - Basic query pattern
  - **DISTINCT**, **ALL**
  - Set combination (**UNION**, **INTERSECT**, **EXCEPT**)
  - Joins
  - Sub-queries (**IN**, **EXISTSSOME**, **ALL**)
  - Aggregate functions
  - **GROUP BY**, **HAVING**
  - View, view updates