

# 403: Algorithms and Data Structures

Prof. Petko Bogdanov

## Introduction

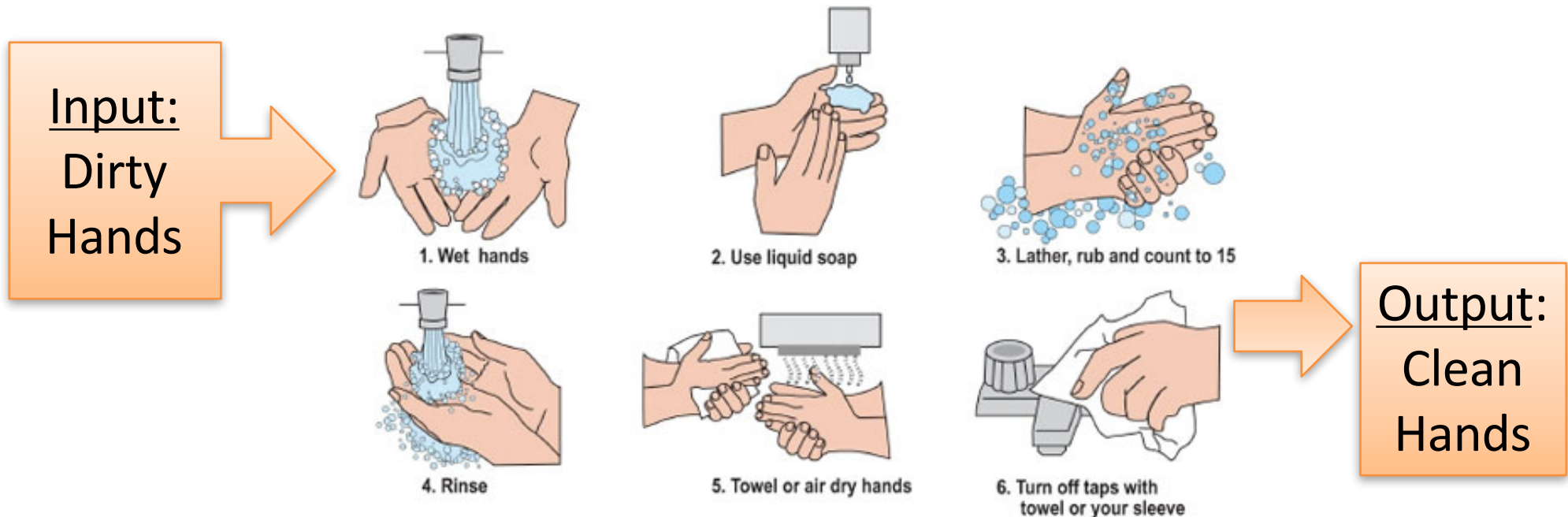
Fall 2016

UAlbany

Computer Science

# What is this course about?

- Design and analysis of algorithms.
- What is an algorithm?
- Algorithm: a well-defined (computational) procedure to solve a problem



# Before we delve into algorithms...

- **Instructor:** Petko Bogdanov,
- For this week only: Prof. Mariya Zheleva  
– mzheleva@albany.edu
- Office Hours: Tue, Thu: 1pm to 3pm, **or by appointment**
- Office: UAB 416 -> Administrative building not on the podium (1215 Western Ave)
- Email: pbogdanov@albany.edu



East Center Drive

Uncommon Grounds  
Coffee & Bagels

1-71 University Place

UAB

Google

# TAs

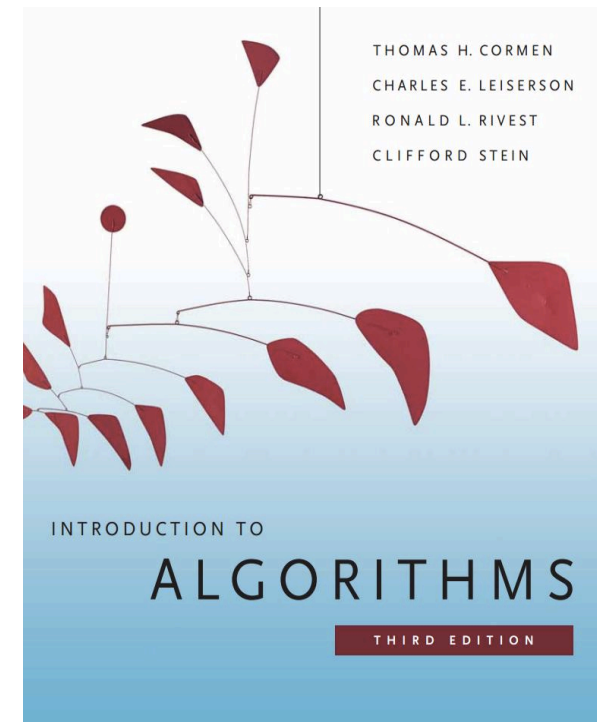
- Ashish Jadhav
  - [ajadhav2@albany.edu](mailto:ajadhav2@albany.edu)
  - Office hours: Tue,Thu 10:30am-11:30am.
- Zumrut Akcam
  - [zakcam@albany.edu](mailto:zakcam@albany.edu)
  - Office hours: Tue 3-4pm and Wed 11am-12pm
- TA office hours rooms TBD (look for an announcement from Blackboard)

# Prerequisites

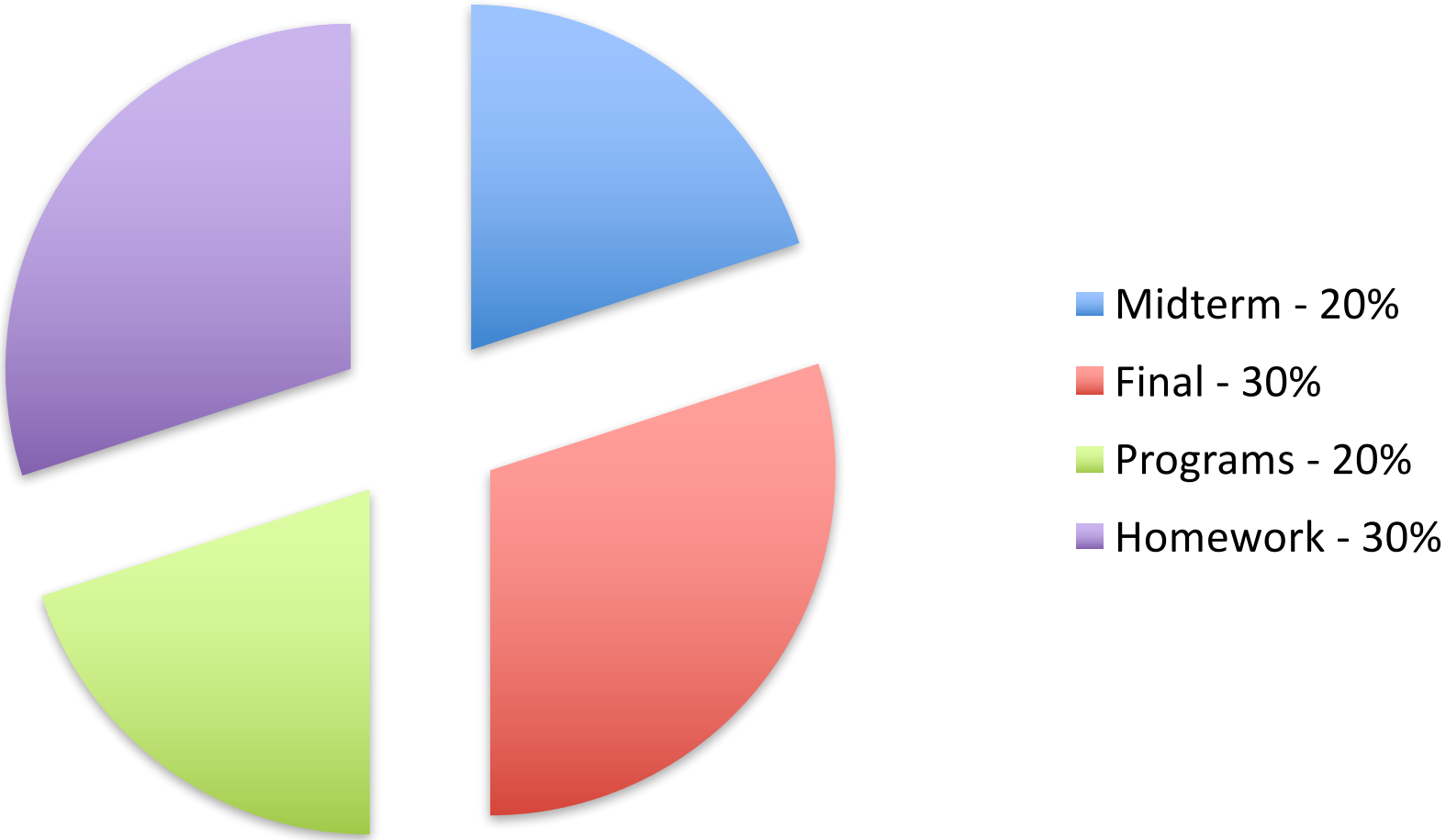
- Prerequisites:
  - CSI 210 (Discrete Structures) and
  - CSI 213 (Data Structures)
- Is there someone who has not taken those classes?
  - Do HW 1 and ...
  - If you get more than 80 points you can stay in the class
  - Else, please, take 210 and 213 first

# Book

- Required Text:
  - Thomas H. Cormen, Charles E. Leiserson, Ronald L. Rivest, Clifford Stein, *“Introduction to Algorithms: Ed. 3”*, MIT Press
- We will mostly follow the book
- Some assignment from the book



# Grading





# Policies: exams

- **Two exams** – open book
  - Midterm – in class on 10/17/2016
  - Final – early December 2016 (date TBD)
- **Make-up exams**
  - Only for major emergencies (valid and verifiable)
  - Contact the instructor as soon as you know you will be missing an exam to re-schedule

# Policies: programming

- **Two programming assignments**
  - Implementation and comparison of algorithms
  - In Java on the department Linux cluster
    - Be sure to check they compile on the cluster before submitting! 0 points for programs that don't compile.
  - Test input/output provided
    - Will not be the same for grading

# Policies: homework

- **Six homework assignments**
  - About a week to complete each
  - Equally-graded
  - Best-5-of-6 policy
    - No make-up homeworks
- **Assignment submission through Blackboard**
  - Programming as a single archive file
  - Homework as a scanned PDF of your solution
    - Only if you are not able to scan before class: you can bring a hard copy

# Policies: cheating

- Cheating on exams
  - You will receive an E-grade and will be reported
- Cheating on assignments
  - **All assignments are to be completed individually**
  - First attempt: all students involved will get 0 points for the assignment
  - Following attempts will result in E-grade; students will be reported for disciplinary action

# Policies: / grades (incomplete)

- Only given for circumstances beyond your control **IF**:
  - Your work is in *good standing* as of the midterm point (10/17/2016)
    - Homework score  $\geq 50\%$
    - Programming completed at least at 50%
    - Midterm grade equivalent of C
  - You are able to provide written documentation to prove you are unable to complete the course

# Policies: attendance

- Not mandatory but will determine your class performance
  - It is your responsibility to find out the material and notes covered in classes you missed
- This is a hybrid slide/board course
  - Always come prepared to take notes

# Reading ahead

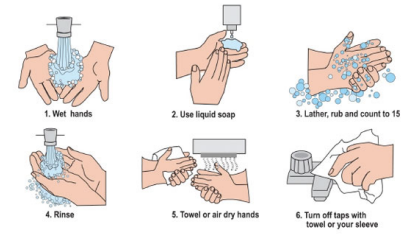
- Read material ahead in the book
  - make more out of the lectures
  - let ideas “sink” longer
  - do not worry if some details unclear
- For this week: Read Chapters 1 and 2

# Homework 1

- Due 09/07/2016 before class
- 4 problems to brush up your Discrete Math
- Need to pass with at least 80% to stay in class
- Start early, so you can consult with the TAs and Prof. Bogdanov.



# Back to our definition



- Algorithm: a well defined (computational) procedure to solve a problem
- To specify the problem we need:
  - Input: What is given?
  - Output: What is needed?

# Example problem specifications

- Finding the maximum
  - Input: A sequence of numbers  $\langle x_1, x_2 \dots x_n \rangle$
  - Output: The largest number among  $x_1, x_2 \dots x_n$
- Sorting in non-decreasing order
  - Input: A sequence of numbers  $\langle x_1, x_2 \dots x_n \rangle$
  - Output: A permutation (reordering) of  $\langle x'_1, x'_2 \dots x'_n \rangle$   
such that  $x'_1 \leq x'_2 \leq \dots \leq x'_n$

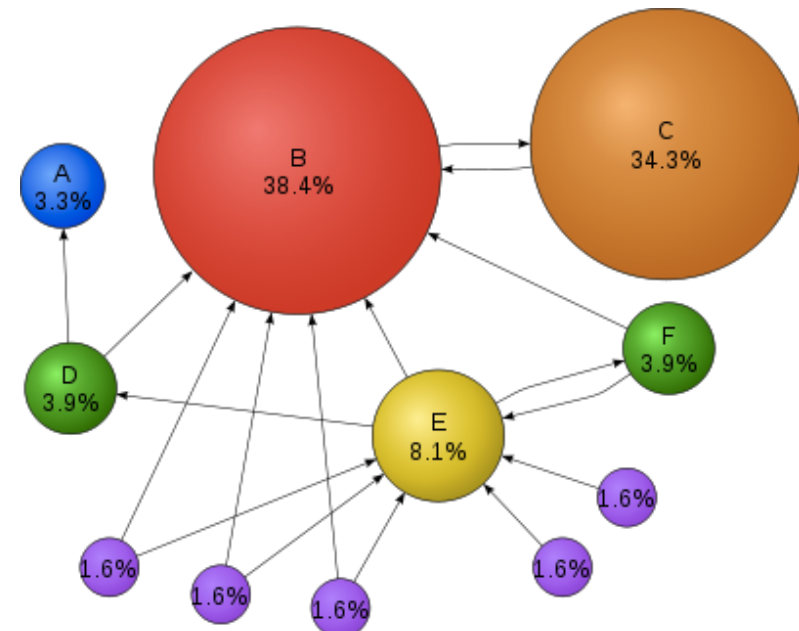
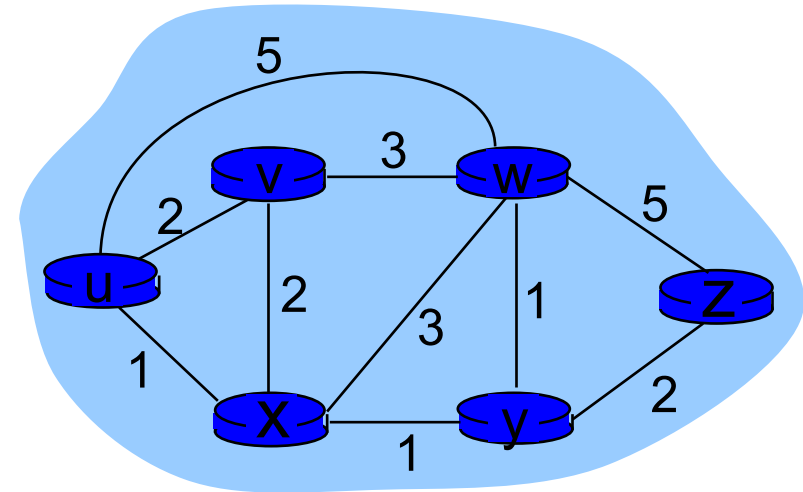
# Instance and correctness

- A sequence such as  $\langle -7, 4, 9, 2, 17, 8 \rangle$  is an **instance** of the problem (i.e. values for all required inputs, satisfying stated constraints)
- A sorting algorithm transforms the sequence  $\langle -7, 4, 9, 2, 17, 8 \rangle$  into  $\langle -7, 2, 4, 8, 9, 17 \rangle$
- An algorithm is **correct** if for every instance it *halts* with the *correct output*.

What kind of problems are  
solved by algorithms?

# The Internet

- Internet routing algorithms
  - Decide how to handle packets through Internet paths
  - You will study Dijkstra's algorithm for finding the shortest path
  
- Internet applications. E.g. Google Search uses PageRank
  - A rough estimate of page importance



# Other examples include

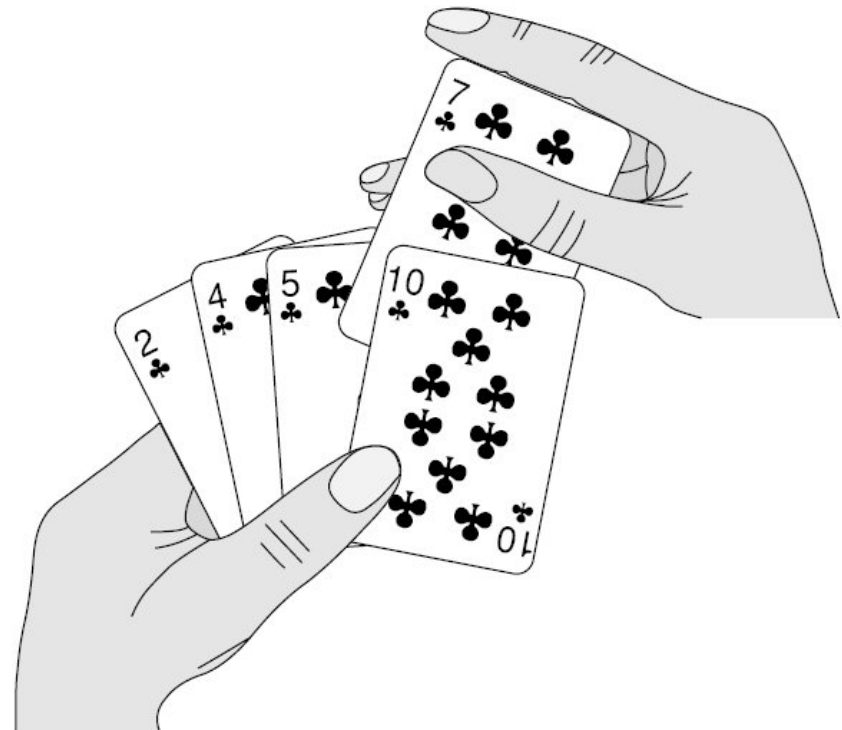
- Electronic commerce
  - Privacy of user information
  - Persistency across multiple server instances
- Manufacturing/commercial/political enterprises
  - Political candidates – how to spend campaign money to maximize the chance of winning
  - Airline companies – assign crews to flights to minimize expenses
  - Internet service provider – where to build infrastructure to minimize delays and maximize Internet bandwidth

# Two common characteristics

- Many candidate solutions
  - Finding one that is “the best” can present quite a challenge
- Practical applications
  - Presents opportunities to narrow the scope → easier to find “the best” solution

# More basics

- Pseudo code and Algorithm Analysis by example
- Insertion sort

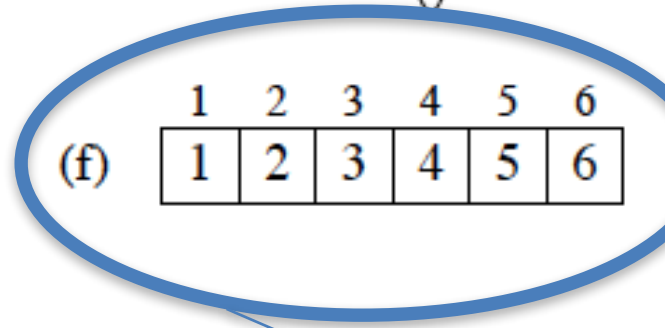
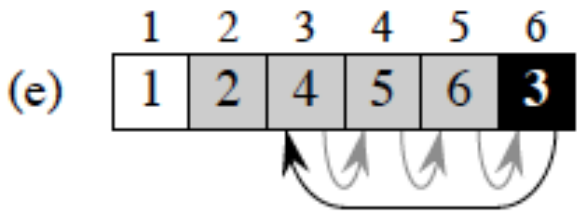
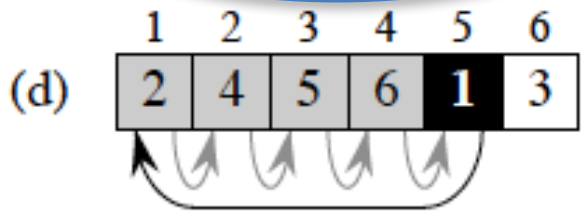
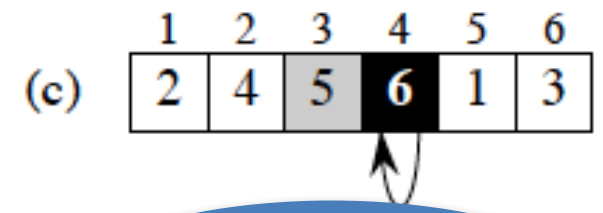
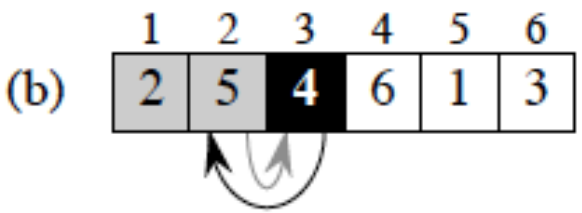
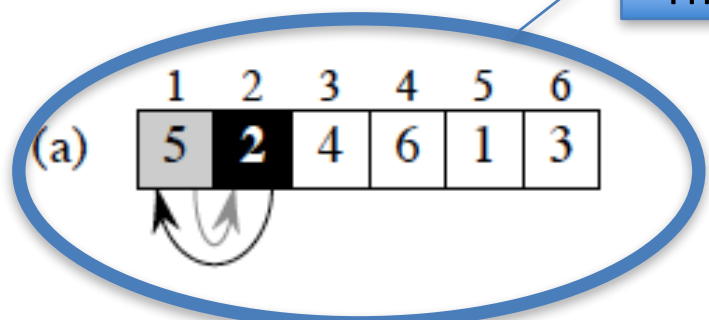




# Example of *insertion sort* on an instance

The algorithm

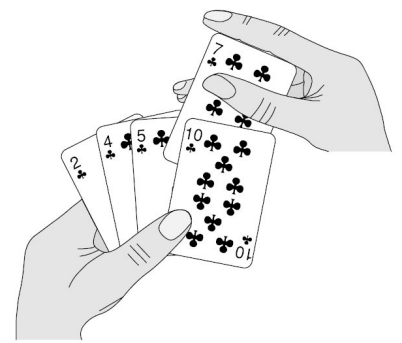
The input



The output

- Which is the algorithm?
- Which is the input?
- Which is the output?
- What is the instance?

<5,2,4,6,1,3>



# Insertion sort (pseudo code)

INSERTION-SORT( $A$ )

```
1  for  $j = 2$  to  $A.length$ 
2       $key = A[j]$ 
3      // Insert  $A[j]$  into the sorted sequence  $A[1 .. j - 1]$ .
4       $i = j - 1$ 
5      while  $i > 0$  and  $A[i] > key$ 
6           $A[i + 1] = A[i]$ 
7           $i = i - 1$ 
8       $A[i + 1] = key$ 
```

This step can be reached when  $i=0$  or if  $A[i] \leq key$ . In both cases key is placed s.t.  $A[1..i]$  is sorted

# What is algorithm analysis?

- Estimating (predicting) the resources needed by an algorithm
  - **Time**: how long would it take
  - **Memory**: how much memory it would require for the used data structures
- Our goal is to find the best algorithm among alternatives in terms of minimum resource requirements.
- How to quantify important algorithm characteristics and omit tedious details?

# Analysis

- Resources are expressed as a function of the size of the input
- The size of the input depends on the problem:
  - *Sorting*: number of elements being sorted  $n$
  - *Integer multiplication*: size of the integers (i.e. number of bits to be multiplied)
- Execution (running) time of an algorithm: the number of primitive operations executed
  - e.g. assignment, comparisons, arithmetic ops. ...
  - typically assumed that operations take const time
  - allows machine independence for the analysis

# Worst case

- When using *running time* we will mean **worst-case** time – the largest running time for any input of a ***fixed size***  $n$ 
  - Note it is also possible to define average (or expected) time assuming we have picked an instance at random from all possible instances, but we will most likely not discuss this here
- In a sense we will be pessimists about our algorithms, always expecting the worst possible input.

# Advantages of being pessimist

- Worst-case time analysis
  - Provides an upper bound on any input, i.e. the algorithm will never take more than this.
  - For some algorithms, worst-case occurs often
  - Helps us provision with redundancy (often a good idea in building software systems)

# Sanity check

- Suppose computers were infinitely-fast and computer memory was free.
  - Would you have any reason to study algorithms?
  - Yes! You still need to demonstrate that your solution **terminates (halts)** and gains the **correct** answer.
- But... memory is not free and CPUs take time to complete a primitive operation.
  - *Efficient* use of resources to gain *correct* answers.

# Announcements



- Next time: Insertion Sort
  - Read through Chapters 1 and 2 from the book
- Homework 1 posted, Due on Sep. 7
- Go over the syllabus on your own and ask questions as necessary